

# Ubiquity Generator Framework: Current Status via a SCIP Application Example

Yuji Shinano  
Zuse Institute Berlin

# Outline

---

- Notes about the UG design and its dynamic load balancing
- Main computational results of ParaSCIP and ParaXpress
  - Solving previously unsolvable instances of MIP
- Some other projects with ParaSCIP
  - A new feature of SCIP is going to be parallelized
  - An application of `ug[SCIP,*]` libraries
- How UG applications were developed
- Future plan
- Concluding remarks



# Outline

---

- Notes about the UG design and its dynamic load balancing
- Main computational results of ParaSCIP and ParaXpress
  - Solving previously unsolvable instances of MIP
- Some other projects with ParaSCIP
  - A new feature of SCIP is going to be parallelized
  - An application of `ug[SCIP,*]` libraries
- How UG applications were developed
- Future plan
- Concluding remarks



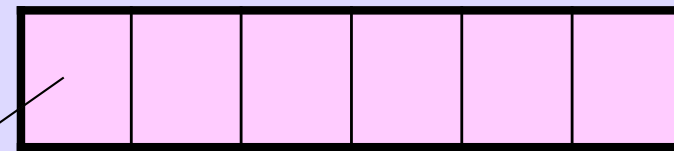
# Initial design of ParaSCIP

## Two Layered Load Balancing

Goal: run with 10,000 cores

### MasterBalancer

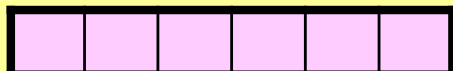
LoadCoordinator statuses



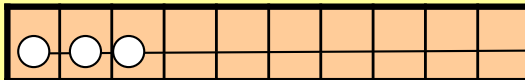
Send a node to the LoadCoordinator

### LoadCoordinator

statuses



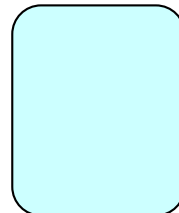
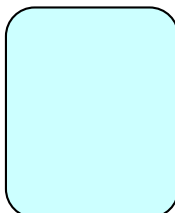
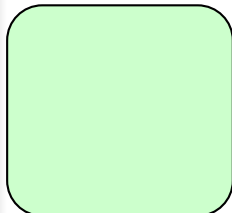
pool



**Solver**

Solver

Solver

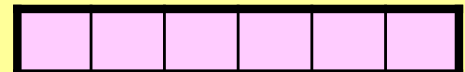


...

...

### LoadCoordinator

statuses



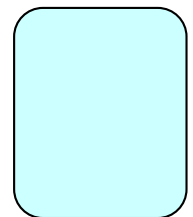
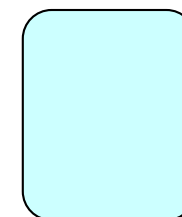
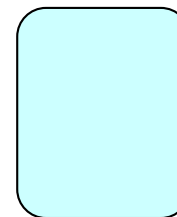
pool



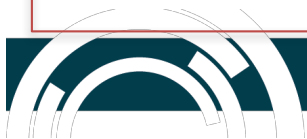
Solver

Solver

Solver



...



# Solving hc9p

Table 1: Statistics for solving hc9p on supercomputers

Run	Computer	Cores	Time (sec.)	Idle (%)	Trans.	Primal bound (Upper bound)	Dual bound (Lower bound)	Gap (%)	Nodes	Open nodes
1	ISM	72	604,796 (317)	< 0.3	738	30,242.0000	29,879.3721	1.21	0	0
						30,242.0000	30,058.9366	0.61	110,012,624	1,257,112
2	ISM	2,304	604,794	< 1.5	979,695	30,242.0000	30,058.7930	0.61	0	15
						30,242.0000	30,102.7556	0.46	3,758,532,600	723,167
3	HLRN III	24,576	86,336	< 1.7	8,811,512	30,242.0000	30,102.6645	0.46	0	35
						30,242.0000	30,116.3592	0.42	2,402,406,311	575,678
4	HLRN III	12,288	43,199	< 1.5	1,709,027	30,242.0000	30,115.3331	0.42	0	3,709
						30,242.0000	30,120.4801	0.40	664,909,985	602,323
5	HLRN III	12,288	118,259	1.5	9,158,920	30,242.0000	30,120.4801	0.40	0	285
						30,242.0000	30,242.0000	0.00	1,677,724,126	0

Supercomputers used:

- ▶ ISM: HPE SGI 8600 with 384 compute nodes, each node has two Intel Xeon Gold 6154 3.0GHz CPUs(18 cores×2) sharing 384GB of memory, , and an Infiniband (Enhanced Hypercube) interconnect
- ▶ HLRN III: Cray XC40 with 1872 compute nodes, each node with two 12-core Intel Xeon Ivy- Bridge/Haswell CPUs sharing 64 GiB of RAM, and with an Aries interconnect



# How open nodes and active solvers evolved (hc9p)

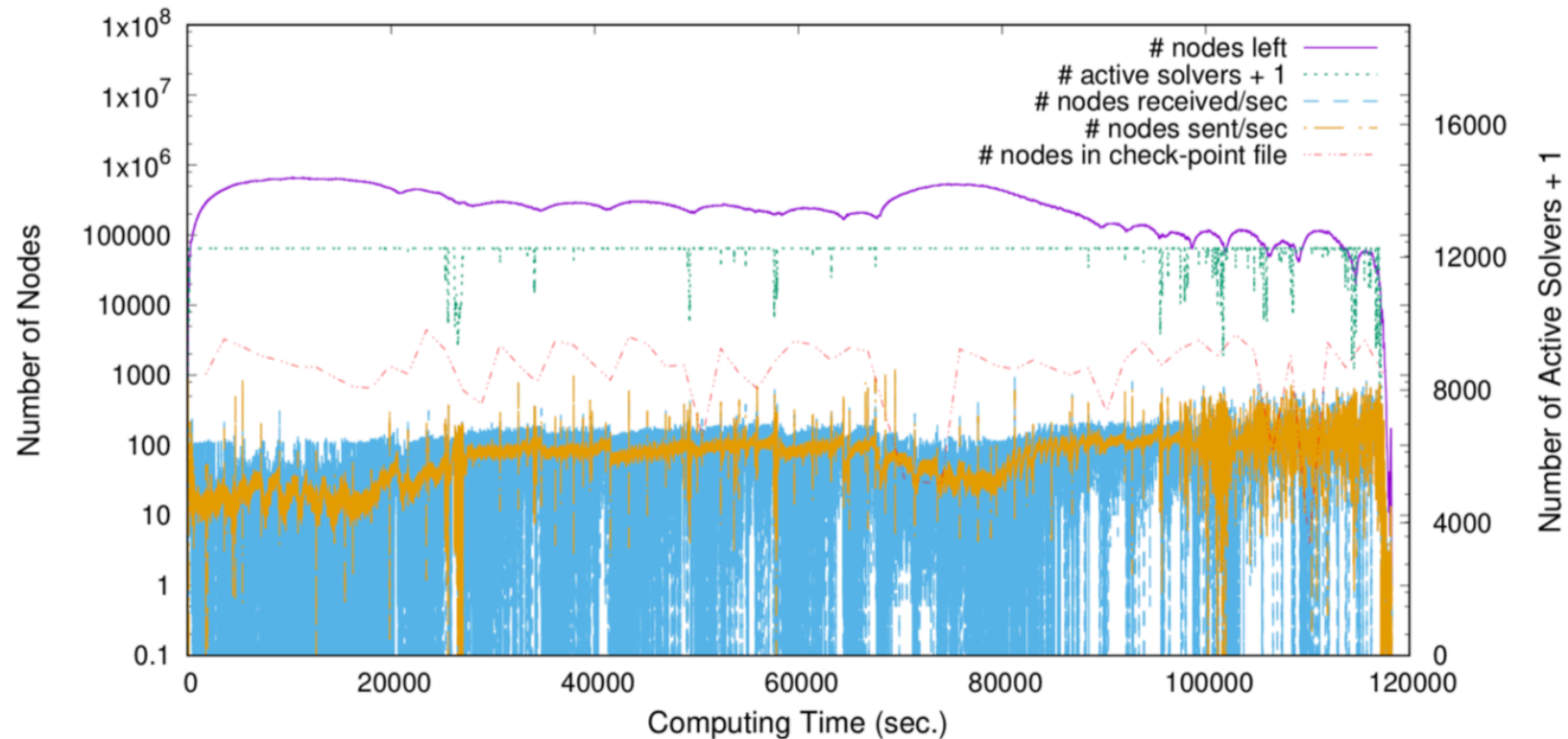


Figure 1: Evolution of computation for solving  $hc9p$  by using 12,288 cores (Run 5)

# Solving hc11p

Table 2: Statistics for solving hc11p on supercomputers

Run	Computer	Cores	Time (sec.)	Idle (%)	Trans.	Primal bound (Upper bound)	Dual bound (Lower bound)	Gap (%)	Nodes	Open nodes
1.1	ISM	72	604,799 (2,558)	< 0.3	71	119,492.0000	117,388.8528	1.79	0	0
						119,297.0000	117,496.5470	1.53	4,314,198	1,109,629
1.2	HLRN III	12,288	43,149 (7,164)	< 0.5	31,304	119,297.0000	117,388.7971	1.63	0	0
						119,297.0000	117,426.2226	1.59	28,491,470	5,433,482
2	HLRN III	43,000	86,354	< 4.9	86,152	119,297.0000	117,426.2226	1.59	0	103
						119,297.0000	117,468.8459	1.56	267,513,609	40,499,188

Supercomputers used:

- ▶ ISM: HPE SGI 8600 with 384 compute nodes, each node has two Intel Xeon Gold 6154 3.0GHz CPUs(18 cores×2) sharing 384GB of memory, , and an Infiniband (Enhanced Hypercube) interconnect
- ▶ HLRN III: Cray XC40 with 1872 compute nodes, each node with two 12-core Intel Xeon Ivy- Bridge/Haswell CPUs sharing 64 GiB of RAM, and with an Aries interconnect



# How open nodes and active solvers evolved (hc11p)

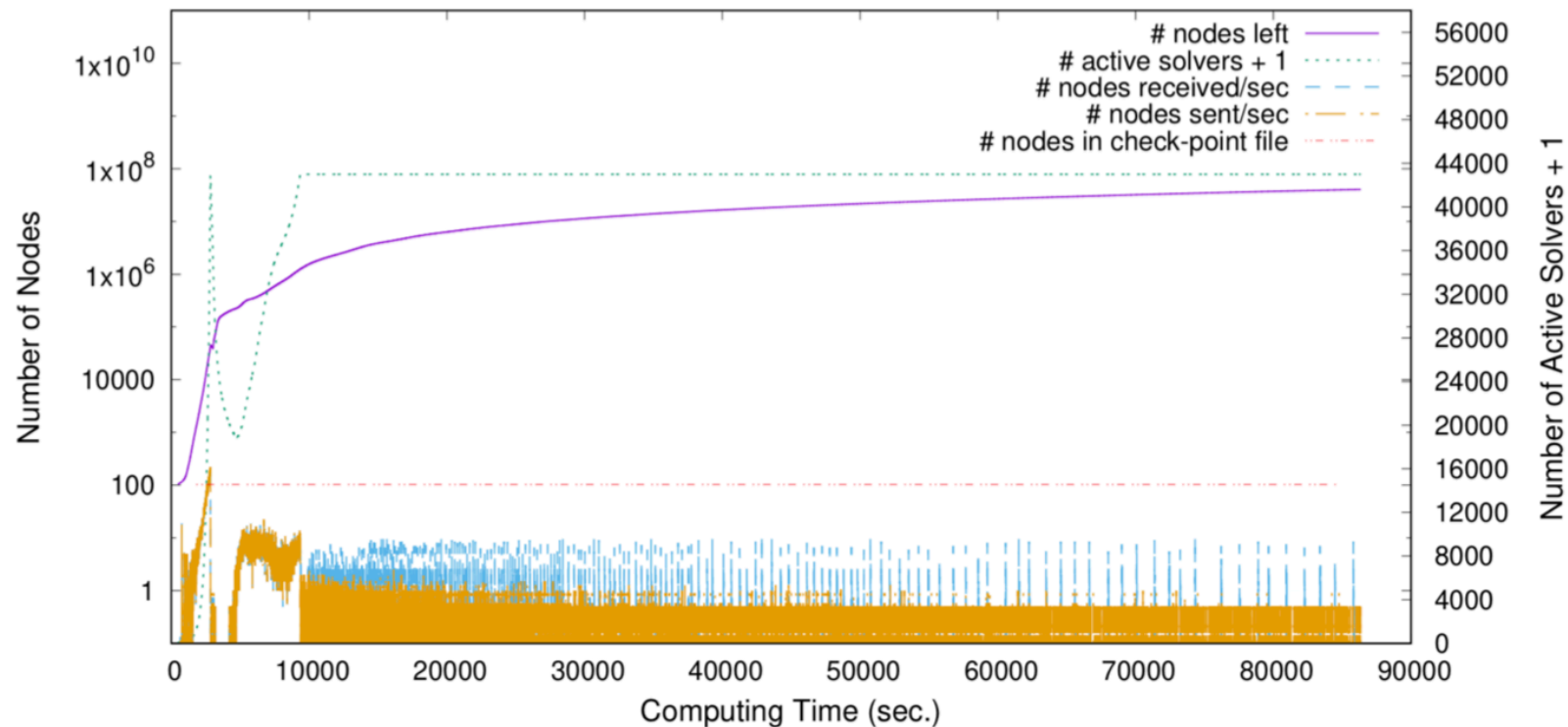
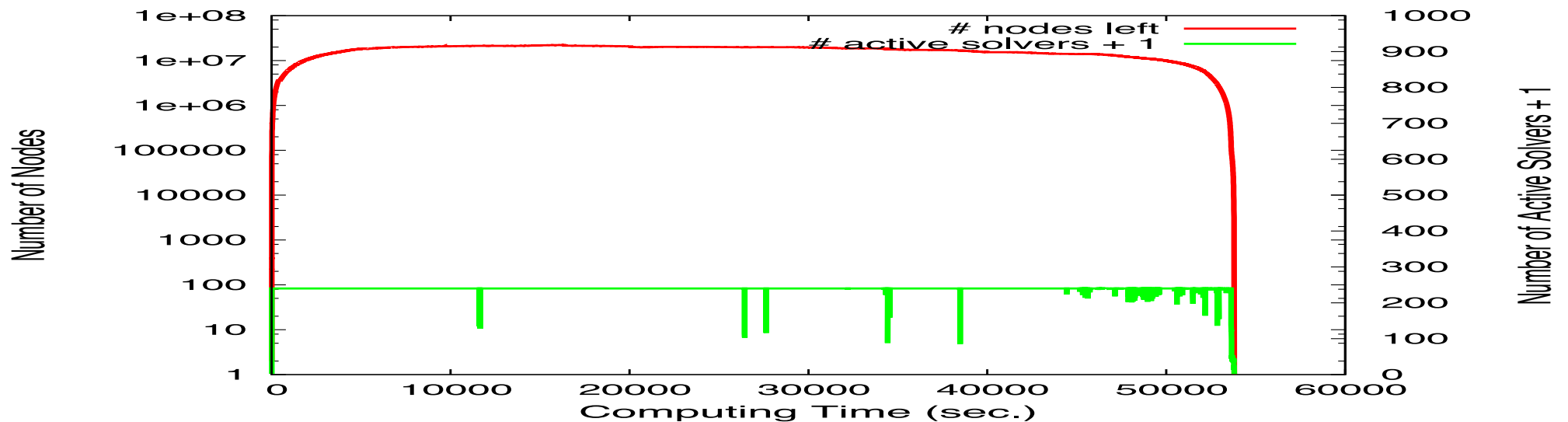
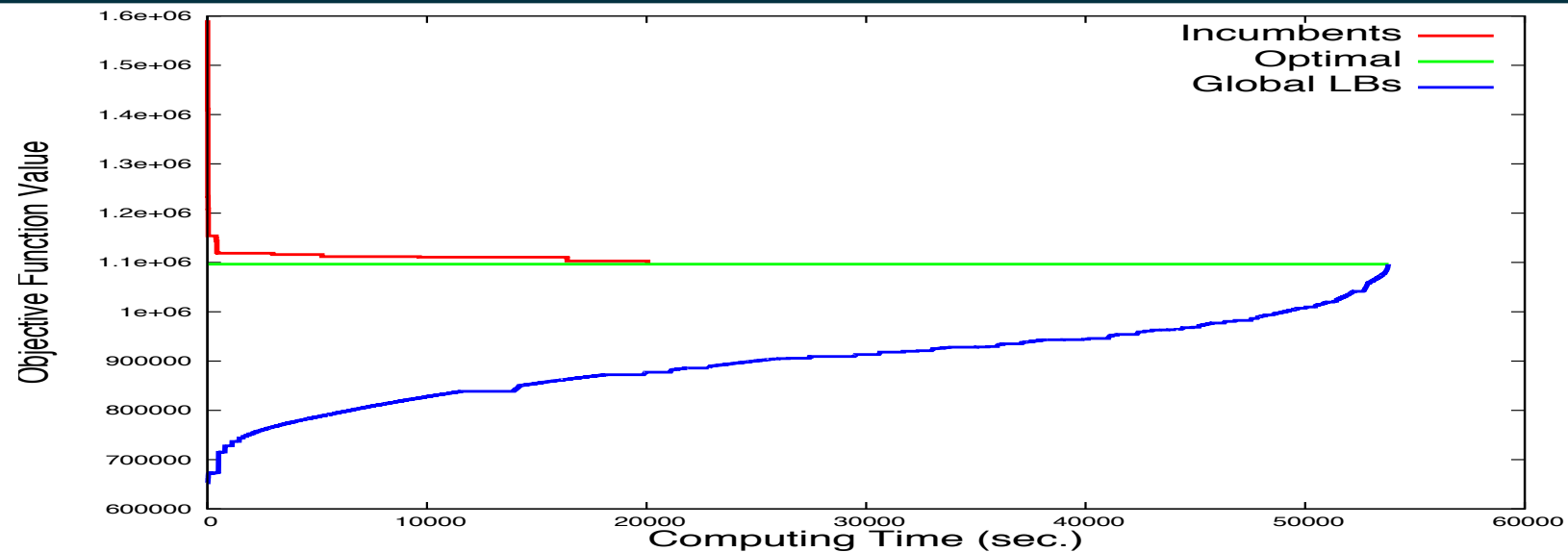


Figure 2: Evolution of computation for solving hc11p by using 43,000 cores (Run 2)

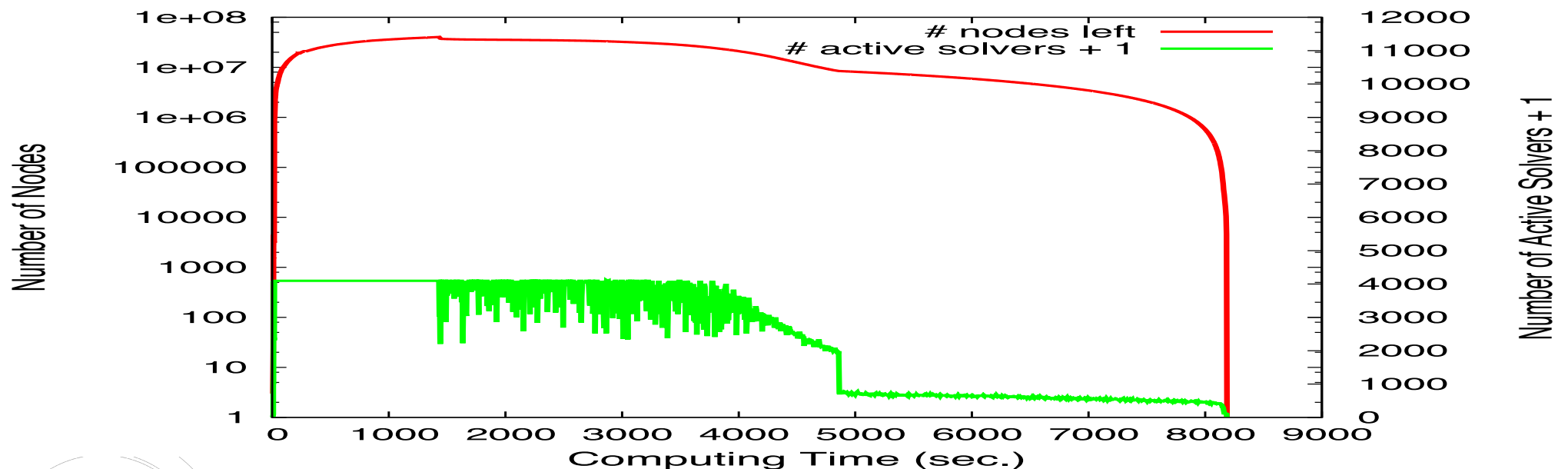
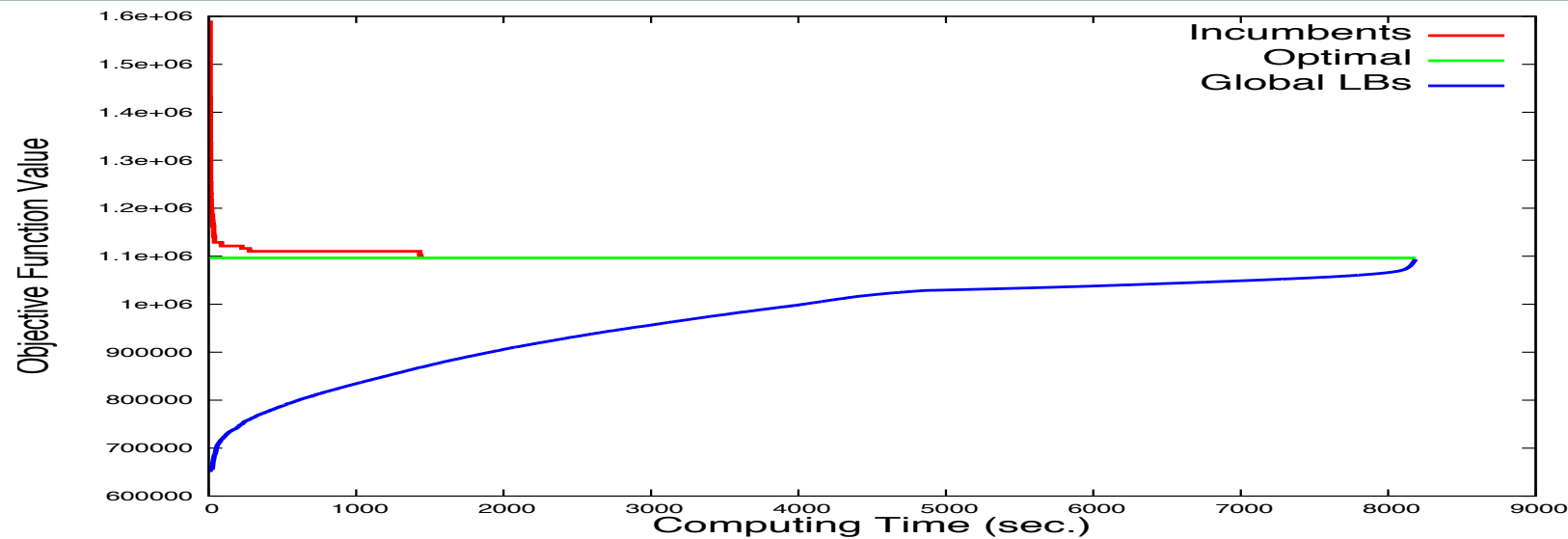




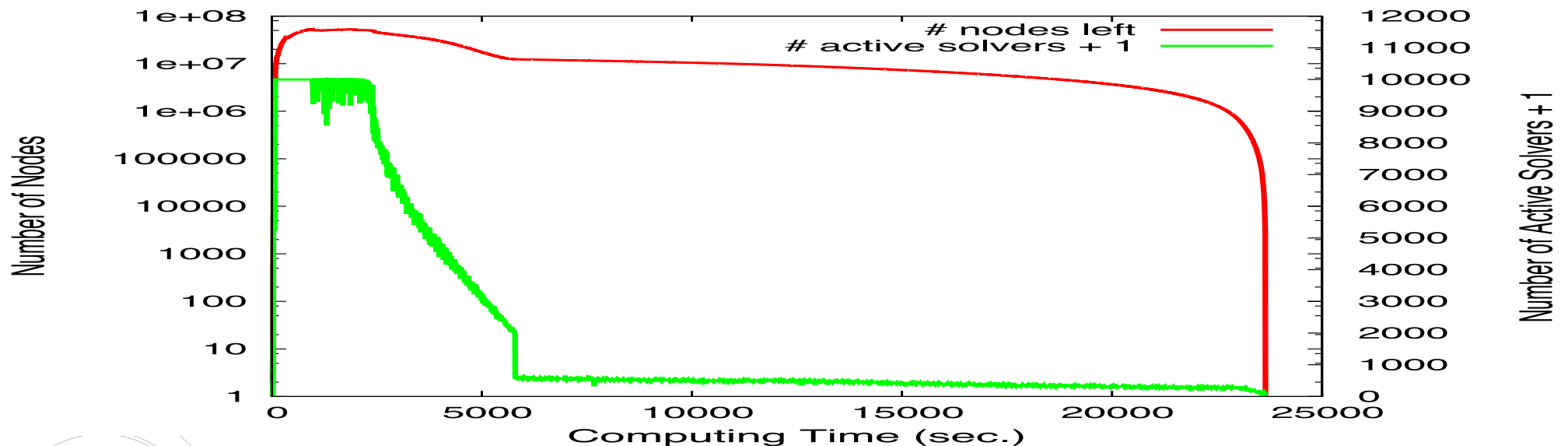
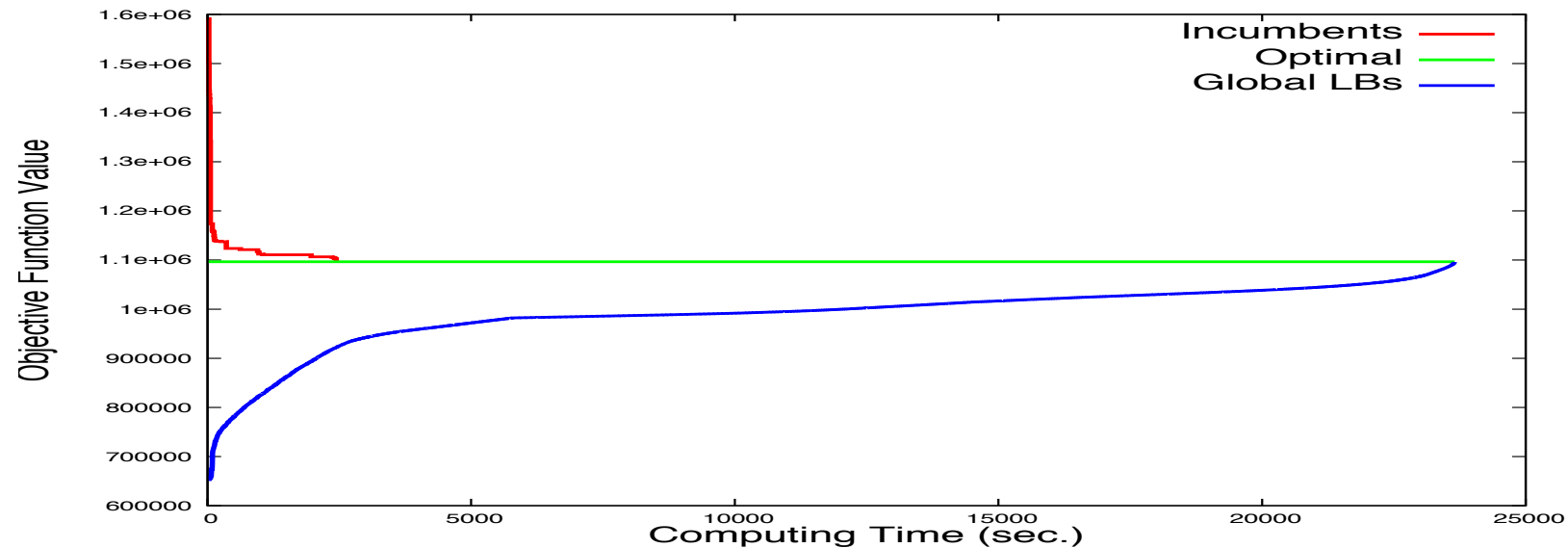
# Runtime behavior: timtab2 – 240 cores



# Runtime behavior: timtab2 – 4096 cores



# Runtime behavior: timtab2 – 10000 cores

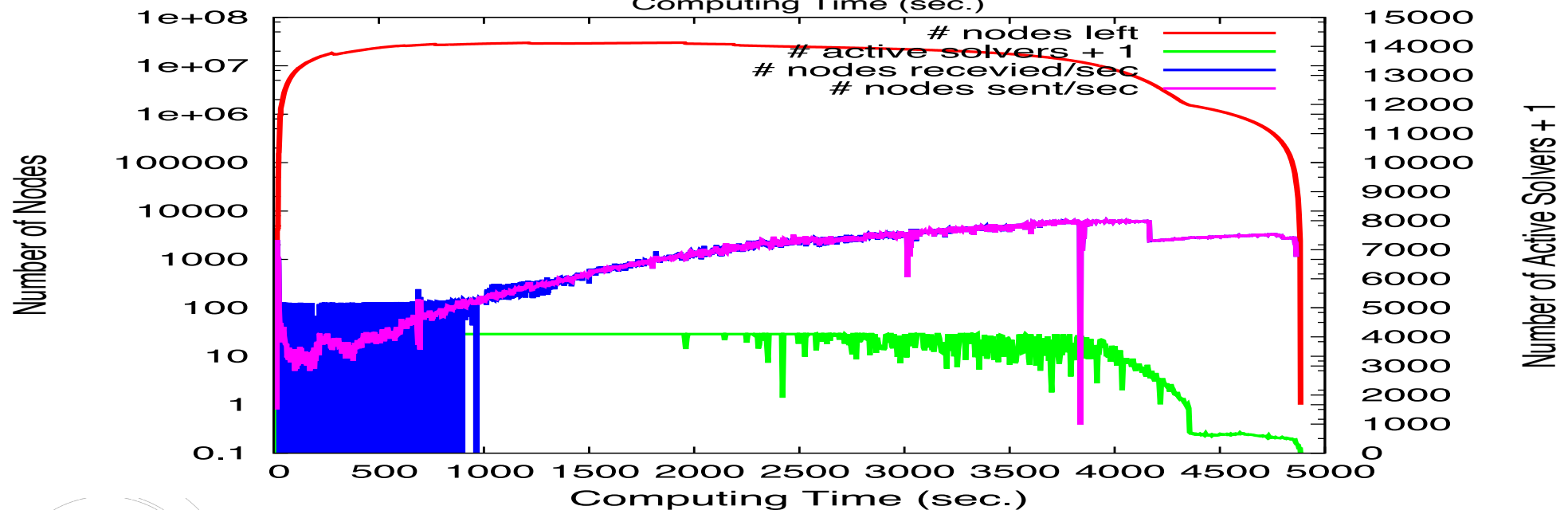
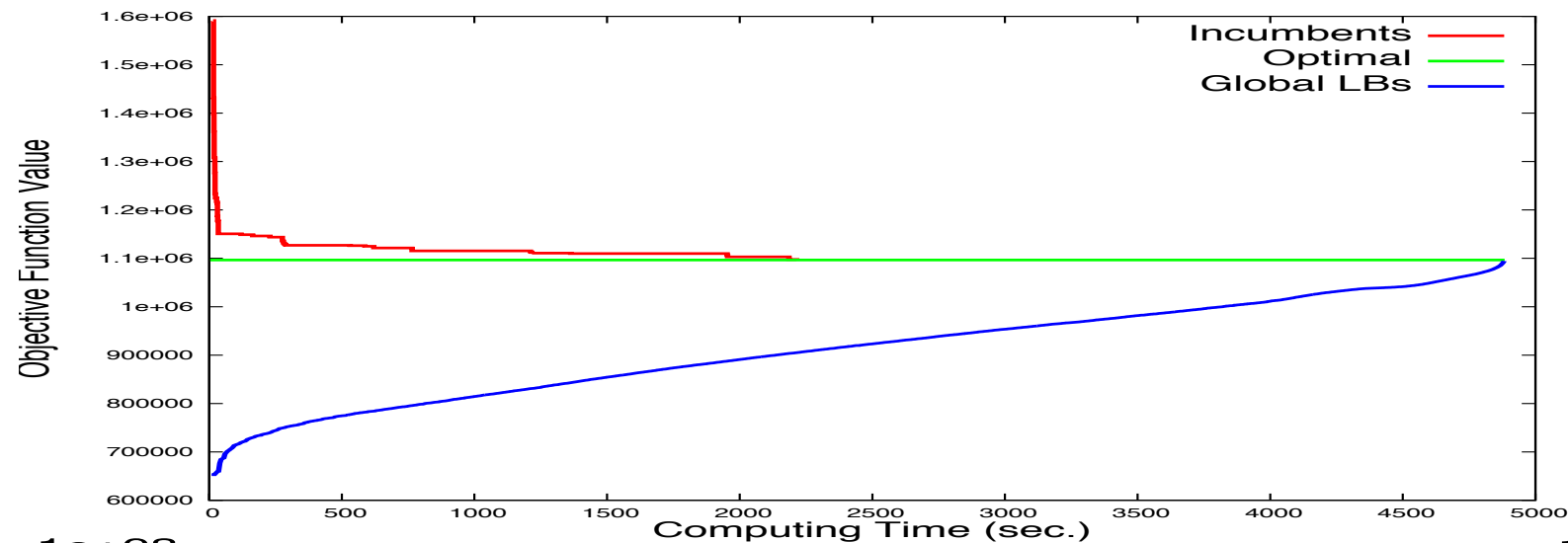


---

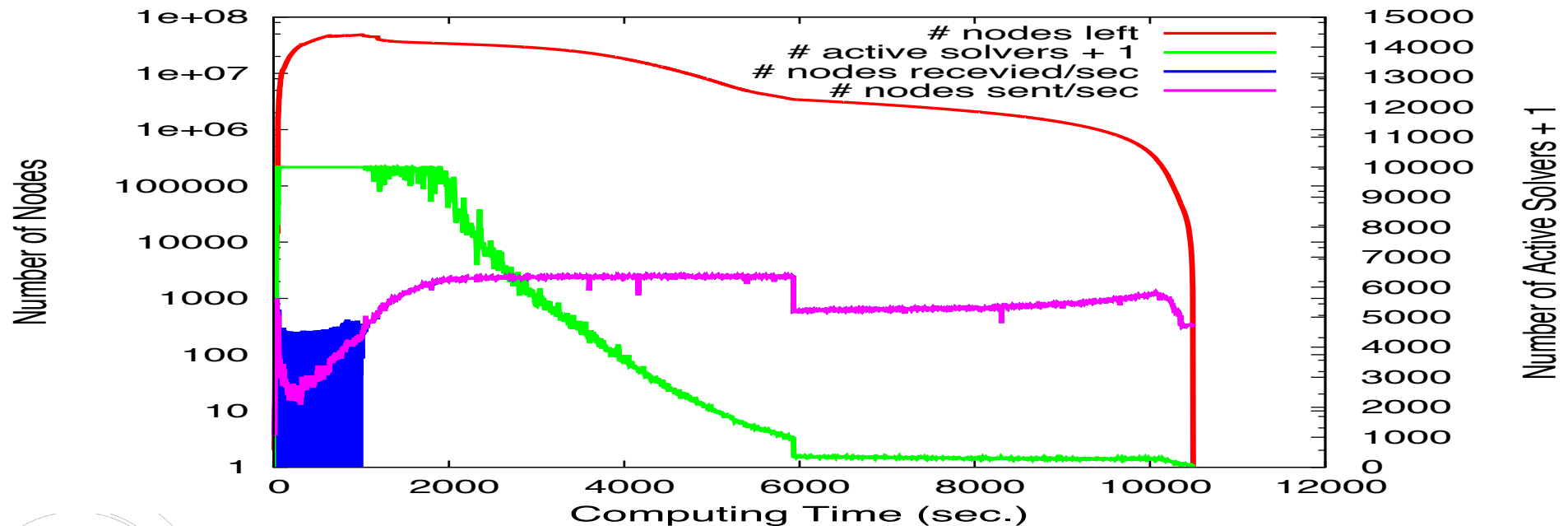
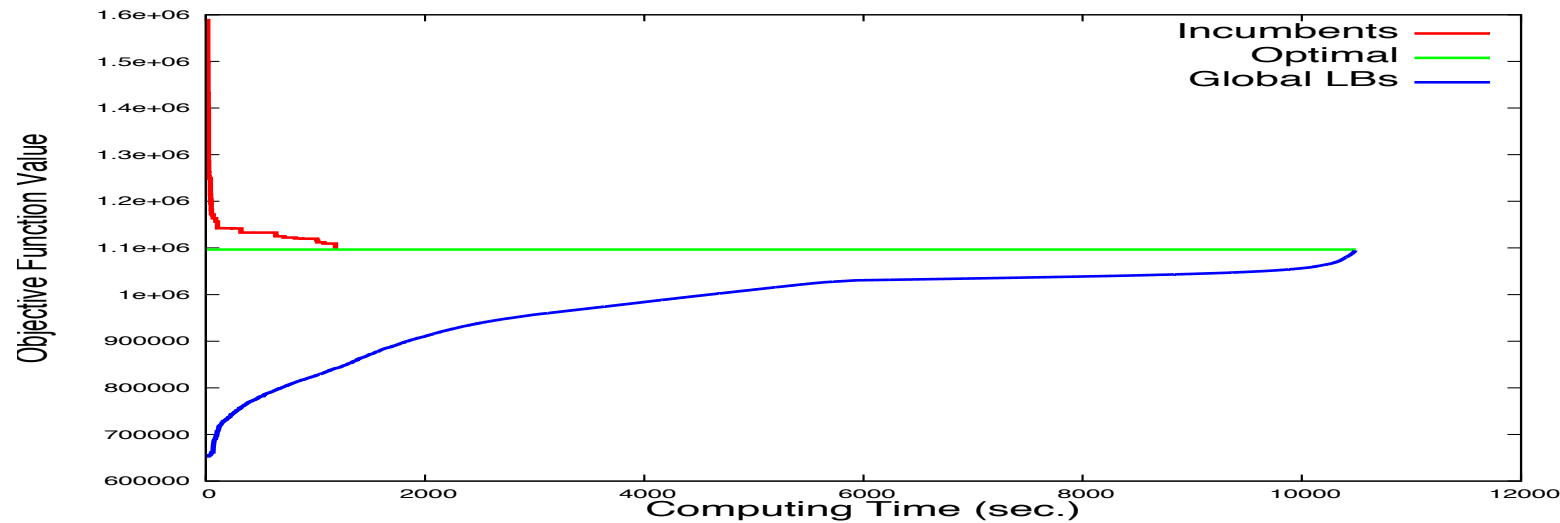
➤ Algorithmic improvements



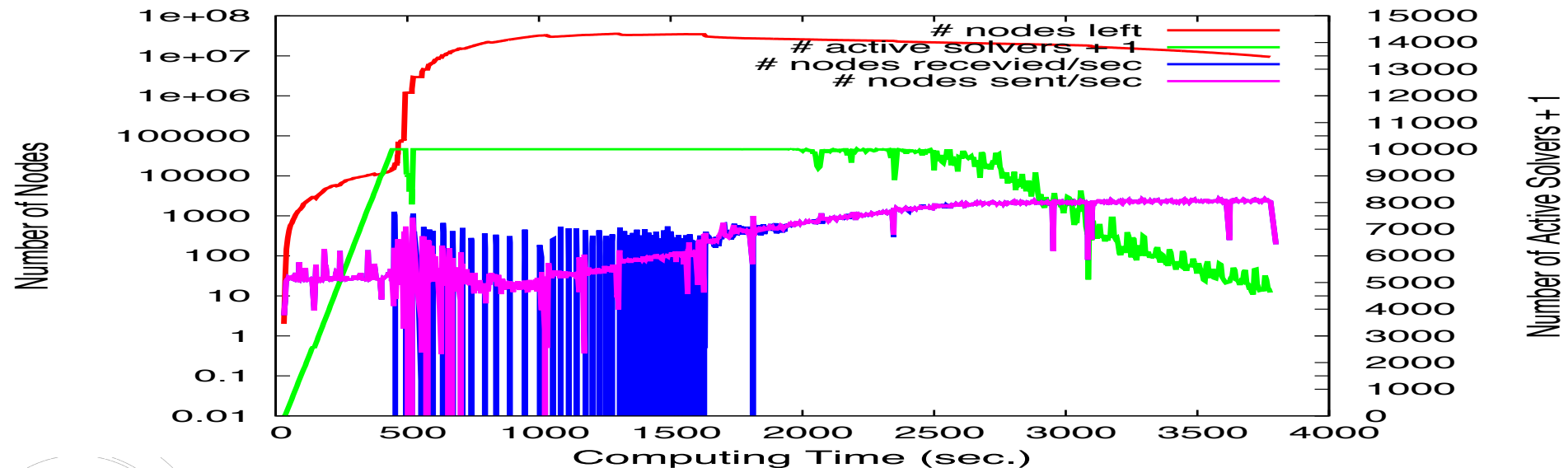
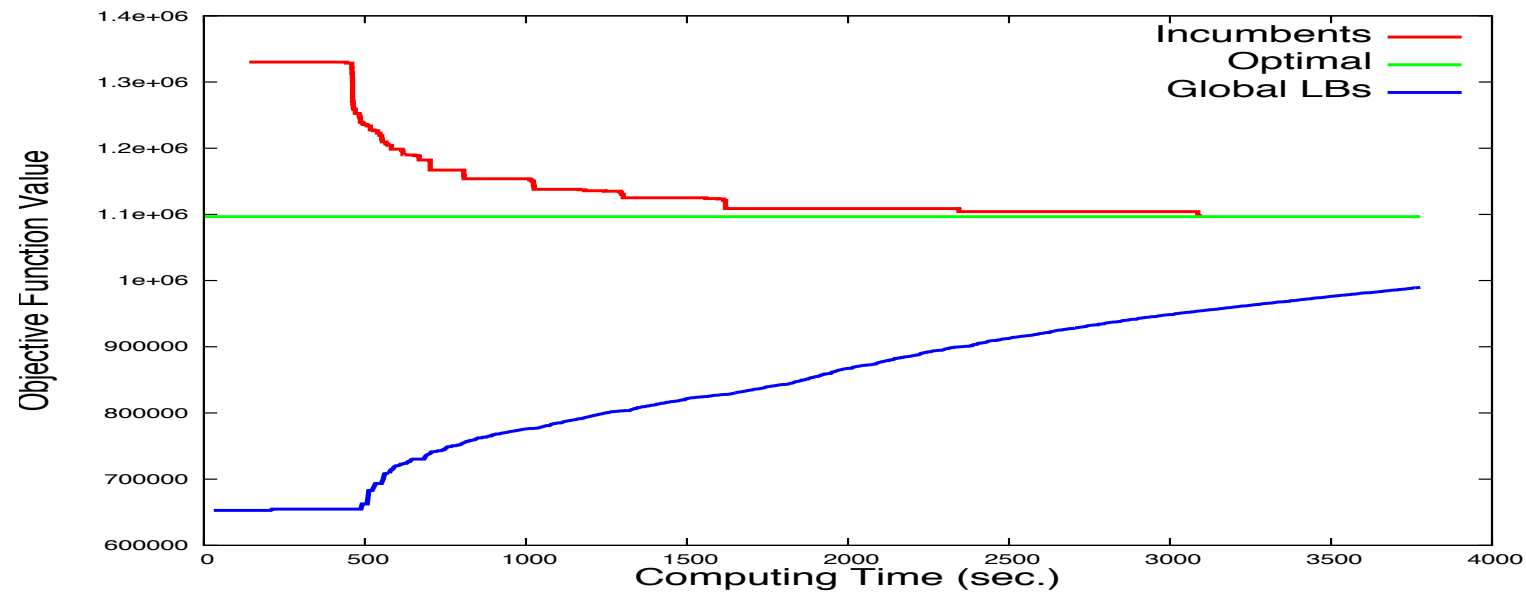
# Runtime behavior: timtab2 (improved) – 4096 cores



# Runtime behavior: timtab2 (improved) – 10000 cores



# Runtime behavior: timtab2 (improved2) – 10000 cores

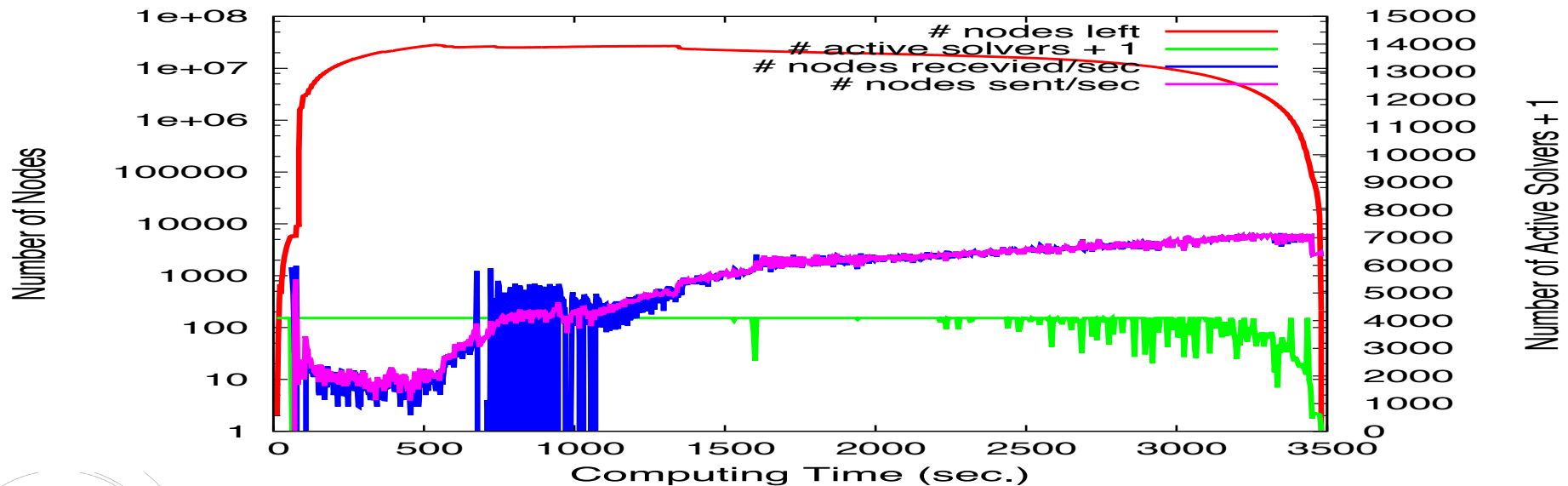
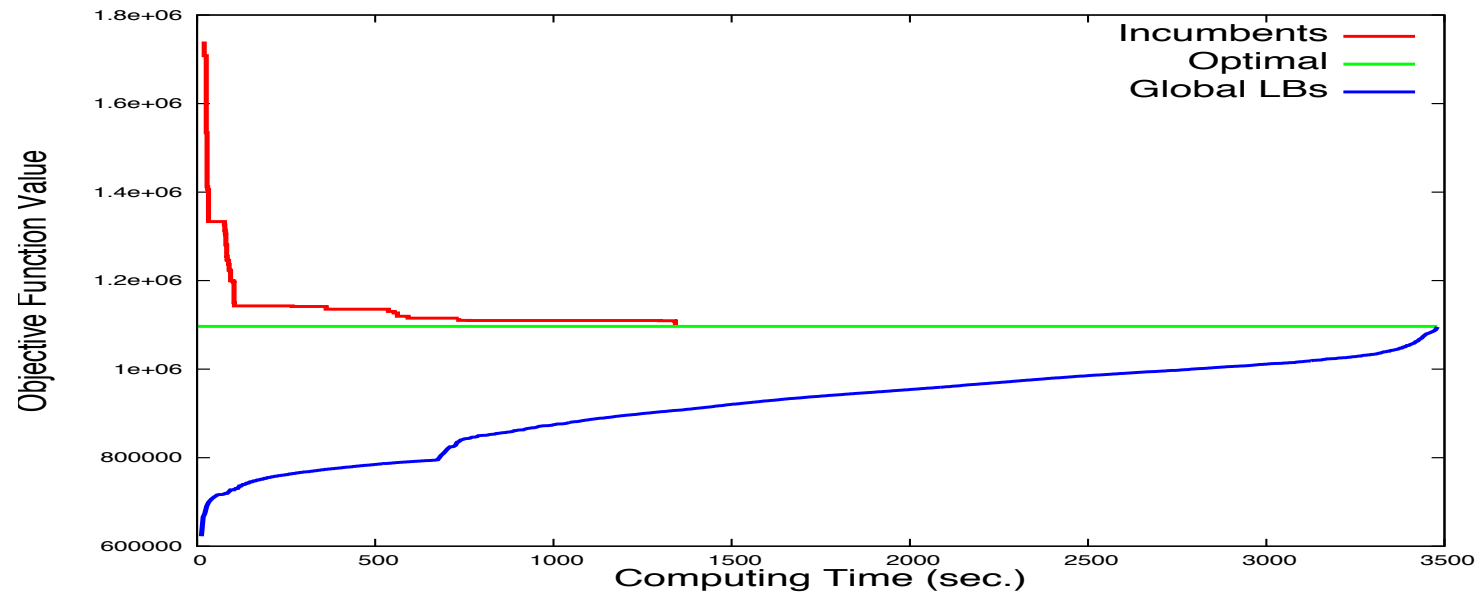


- 
- Improvements of transfer methods  
(with racing ramp-ups)

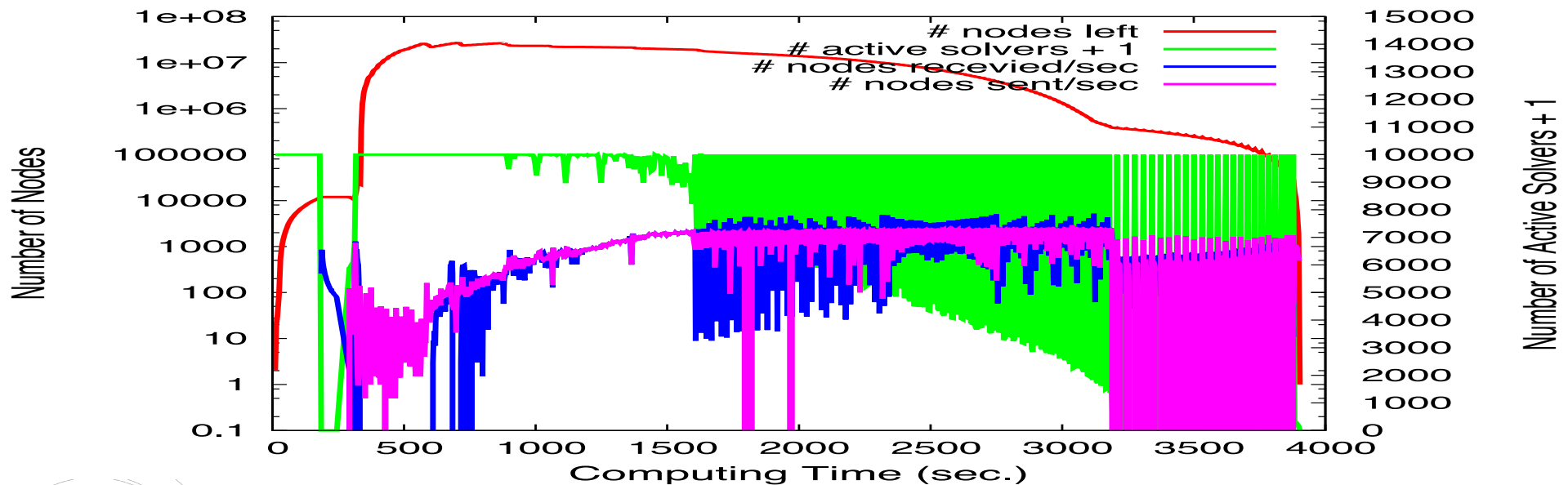
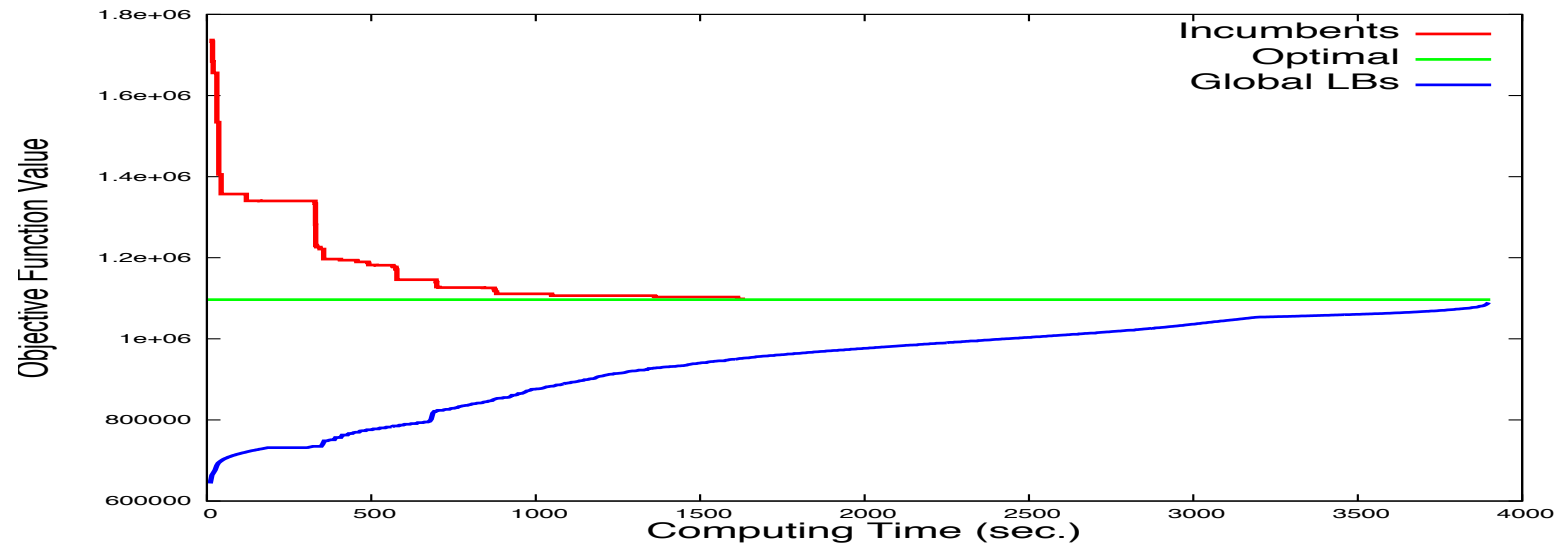




# Runtime behavior: timtab2 (racing) – 4096 cores



# Runtime behavior: timtab2 (racing) – 10000 cores



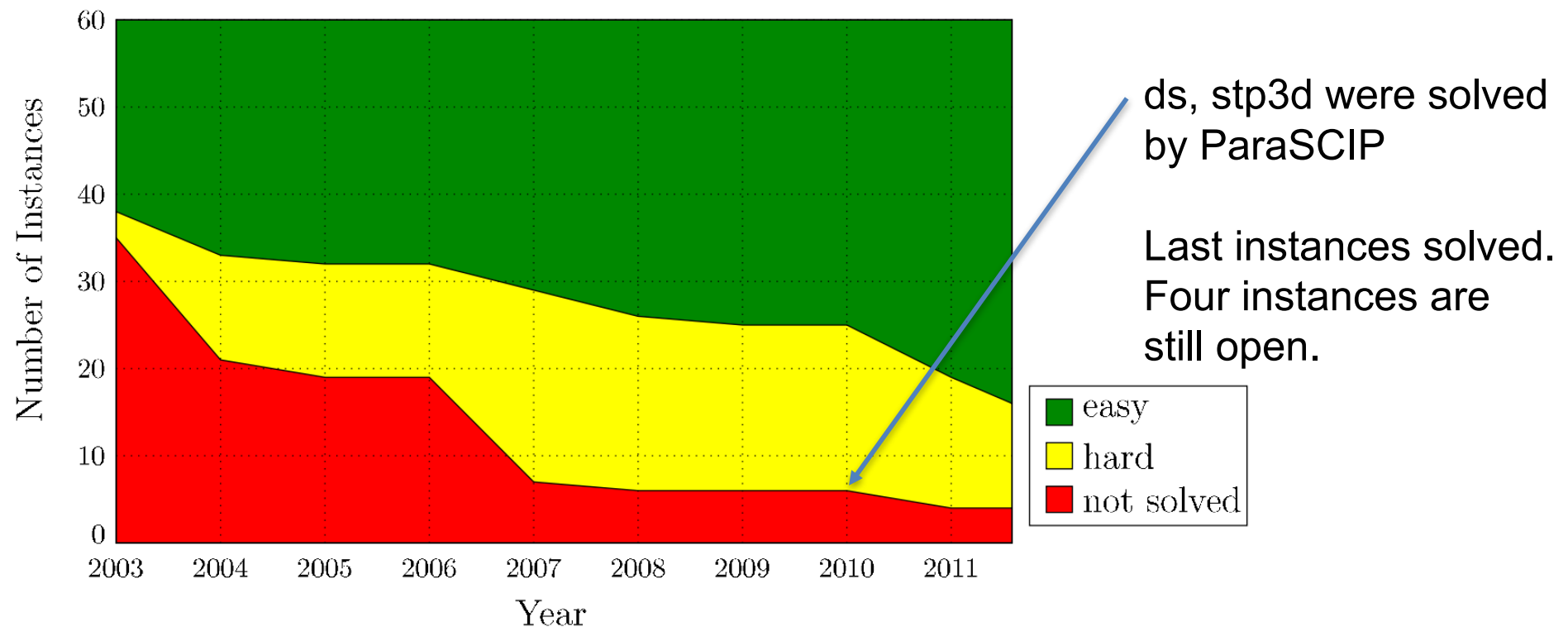
# Outline

---

- Notes about the UG design and its dynamic load balancing
- Main computational results of ParaSCIP and ParaXpress
  - Solving previously unsolvable instances of MIP
- Some other projects with ParaSCIP
  - A new feature of SCIP is going to be parallelized
  - An application of ug[SCIP,\*] libraries
- How UG applications were developed
- Future plan
- Concluding remarks



# Solving open instances from MIPLIB2003



6 instances have not been solved 7 Years

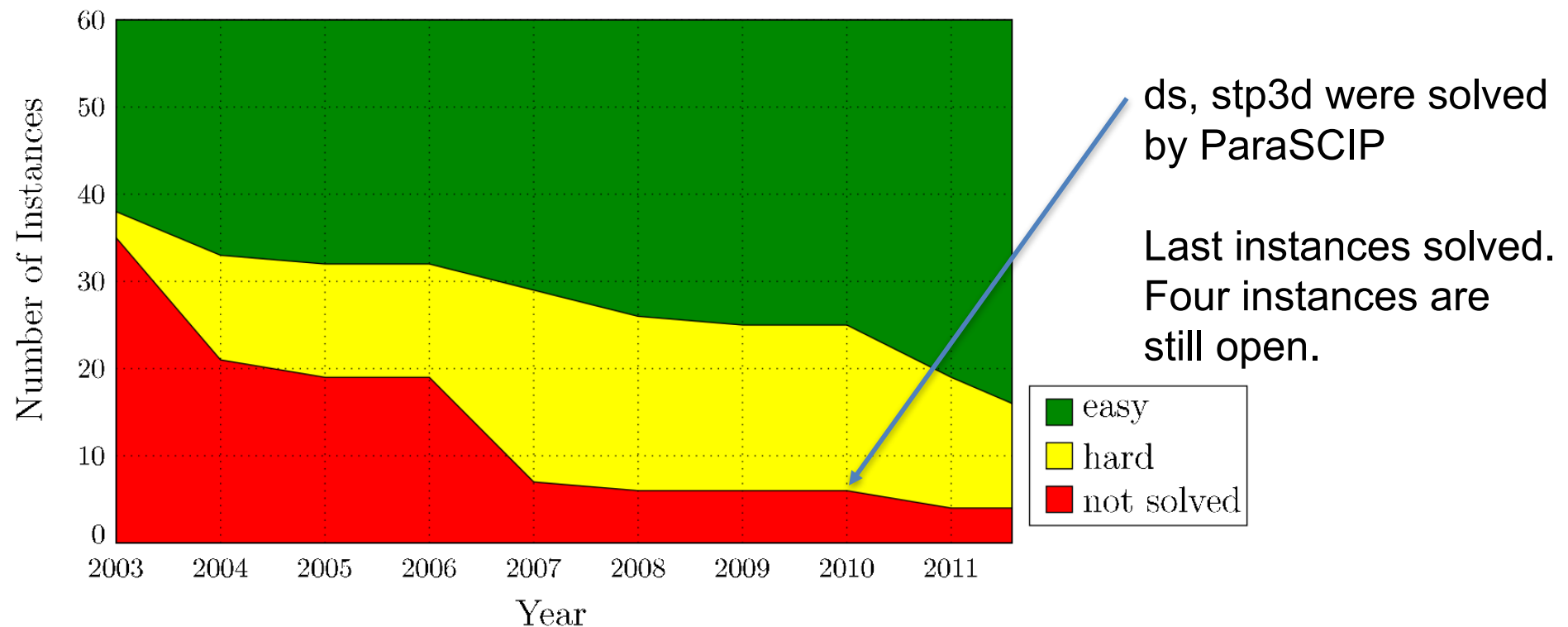
- ds – 656 constraints, 67,732 binary variables
- stp3d – 159,488 constraints, 204,880 binary variables

After applying SCIP presolving 9 times:

88,388 constraints, 123,637 binary variables



# Solving open instances from MIPLIB2003

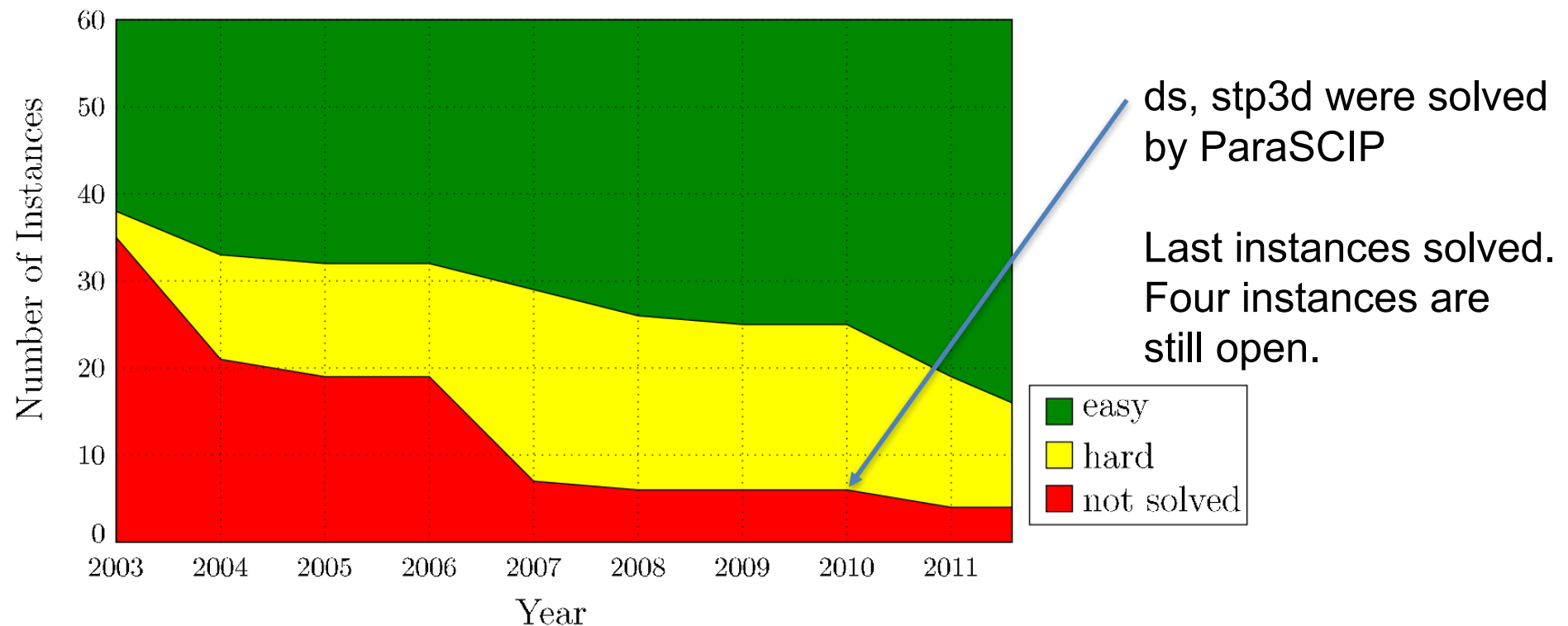


ds: Set partitioning problem originating from public transport service planning.

This instance was solved by a first implementation of ParaSCIP using up to 2048 cores of HLRN-II (<http://www.hlrn.de>). ParaSCIP, mainly developed by Yuji Shinano, is an extension of SCIP and realizes a parallelization on a distributed memory computing environment. For being able to interrupt and warmstart the computations, ParaSCIP has a checkpoint mechanism. Therefore, selected subproblems are stored as warm start information, which allows to virtually run ParaSCIP, although the HLRN-II environment imposes a time limit of 48 hours per run. It took approximately 86 hours to solve this instance.



# Solving open instances from MIPLIB2003



stp3d: 3D Steiner Tree packing problem (VLSI routing problem in a multi-layer grid graph)

This instance was solved by a first implementation of ParaSCIP using up to 2048 cores of HLRN-II(<http://www.hlrn.de>). ParaSCIP, mainly developed by Yuji Shinano, is an extension of SCIP and realizes a parallelization on a distributed memory computing environment. For being able to interrupt and warmstart the computations, ParaSCIP has a checkpoint mechanism. Therefore, selected subproblems are stored as warm start information, which allows to virtually run ParaSCIP, although the HLRN-II environment imposes a time limit of 48 hours per run. The problem was presolved several times with SCIP presolving techniques. After that, it took approximately 114 hours to solve this instance.



# Restarted runs vs. single run

- Which run was efficient from a view point of solving ds?
  - restarted runs: [STAGE 1 (old)]
  - approximated comp. time: 86 [h]
  - real resource requested time: 96 [h]
  - cumulative time: 181,248 [h]
- single run: [STAGE 2(new)]
  - comp. time: 76[h]
  - cumulative time:  $4096 \times 76 = 311,296$  [h]

HLRN II supercomputer (SGI Altix ICE 8200EX)  
maximum cores for a job: 4096

Answer is “restarted runs”

run	cores	time[h]
1	512	4
2	1024	5
3	2048	10
4	2048	7
5	2048	4
6	2048	5
7	2048	4
8	2048	5
9	2048	5
10	2048	4
11	2048	4
12	2048	5
13	2048	10
14	2048	4
15	2048	4
16	2048	12
17	2048	4
Summary		96
Accumulated time		181248



# FiberSCIP and its paper

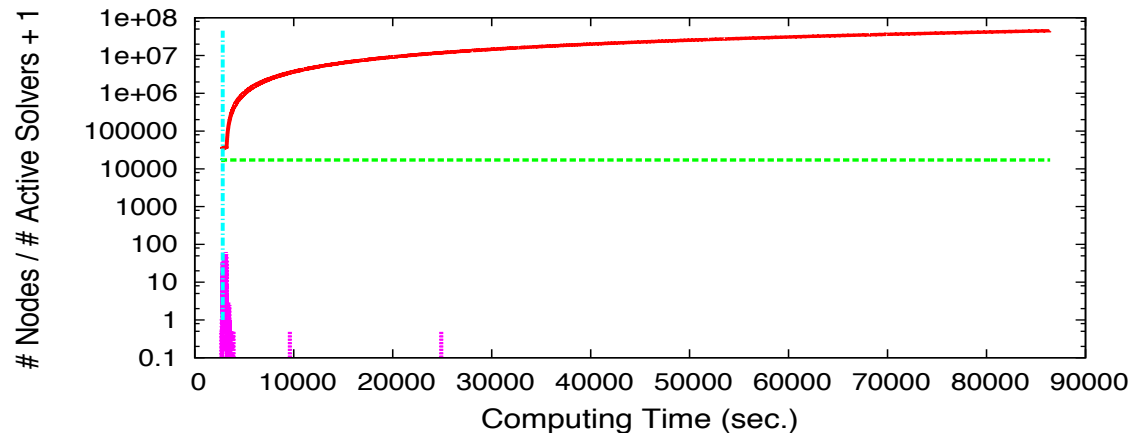
---

- Preliminary results were presented in MIPLIB2010 paper (Started working in 2010)
- Computational results had been checked two years
  - 87 instances of many different settings with five repeated runs
  - All numbers in the [supplement](#) were checked carefully
- Submitted the paper in 2013
- Publish the paper in 2018

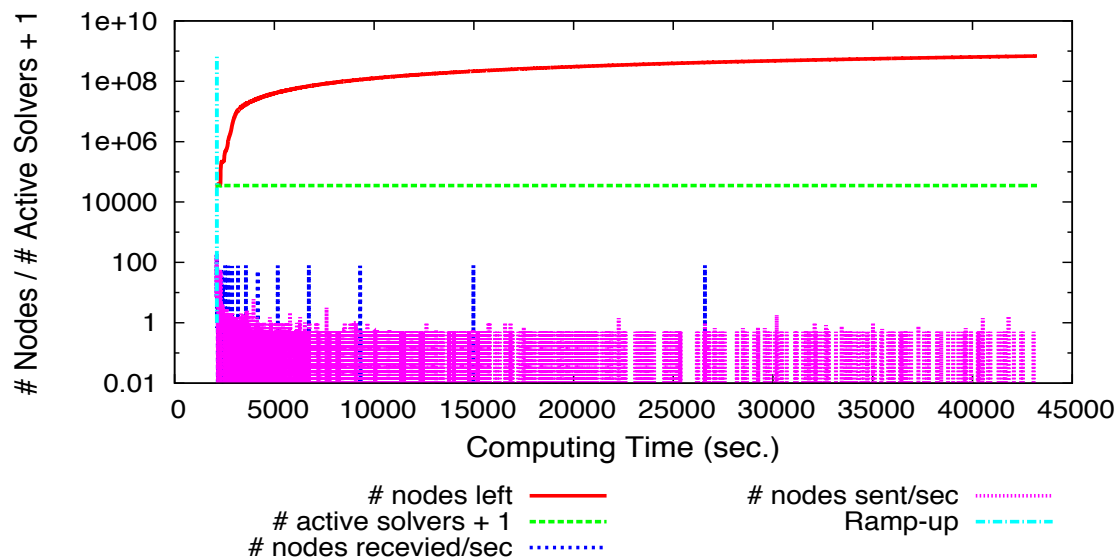




# What does scale up do for a hard instance? (dano3mip)



on HLRN III  
by using 17,088 cores  
starting from 33,332  
branch-and-bound nodes



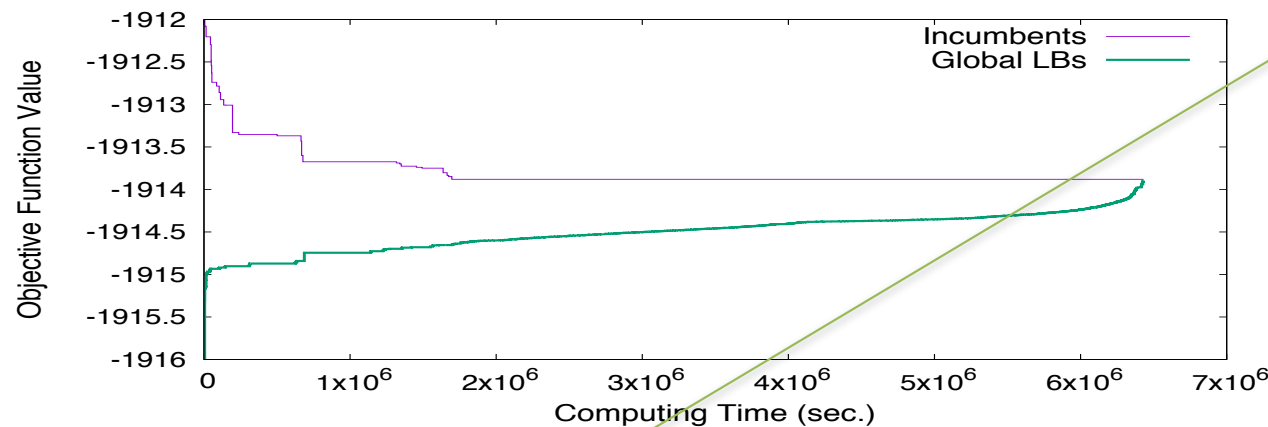
on Titan  
by using 35,200 cores  
starting from 33,481  
branch-and-bound nodes

**Large scale can  
break down  
problem into  
solvable sub-MIPs**



# The biggest and the longest computation

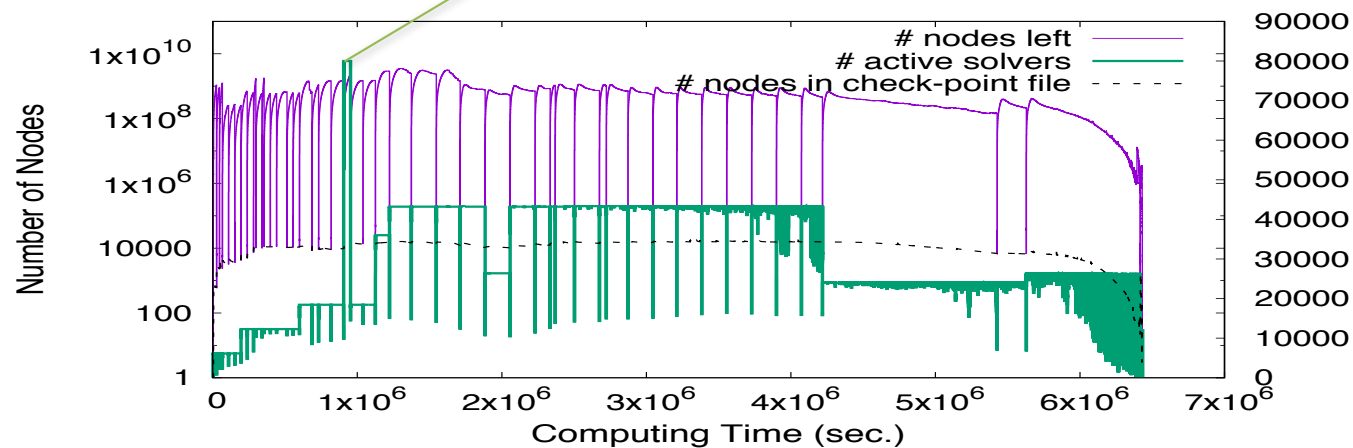
Solving rmine10: 48 restarted runs with 6,144 to 80,000 cores



How upper and lower bounds evolved

Titan with 80,000 cores  
The others: HLRN III

It took about 75 days and  
**5,660 years** of CPU core  
hours!



How open nodes and active solvers evolved

Number of Active Solvers + 1

UG can  
handle up  
to 80,000  
MPI process



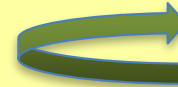
# ParaXpress: Need to cooperate with Xpress developers

$$\text{Original } \min\{c^T x : Ax \leq b, l \leq x \leq u, \text{ for all } x_j \in \mathbb{Z}^n, j \in I\}$$

$$\min\{c'^T x' : A'x' \leq b', l' \leq x' \leq u', \text{ for all } x'_j \in \mathbb{Z}^{n'}, j \in I'\}$$

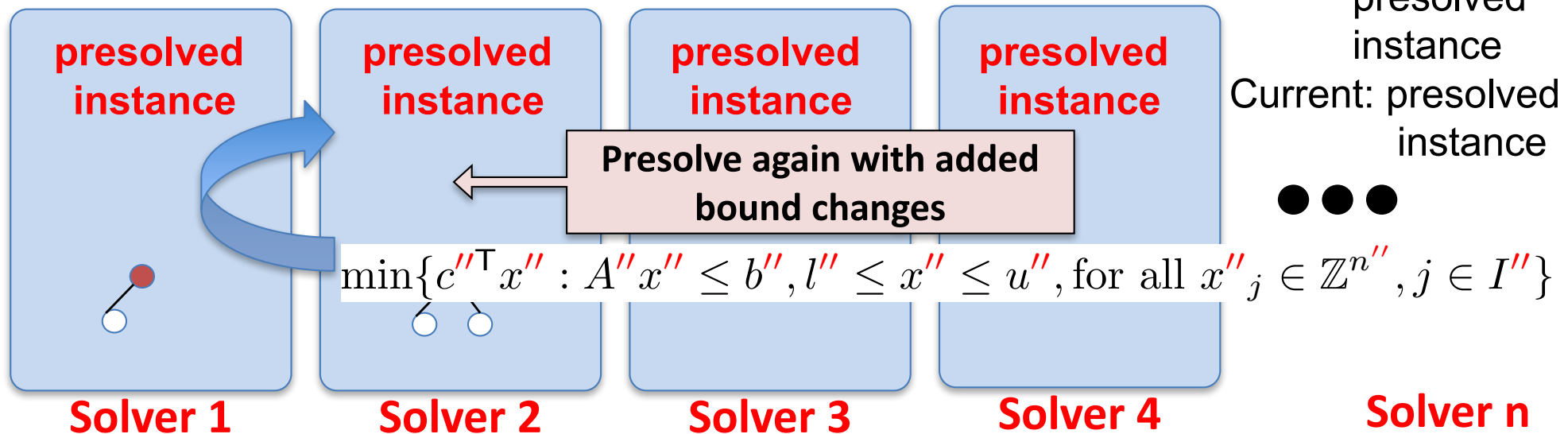
## LoadCoordinator

waiting:  $(l'_i, u'_i)$   
running:



Base solver  
I/O, **presolve**

First: restricted  
presolved  
instance  
Current: presolved  
instance



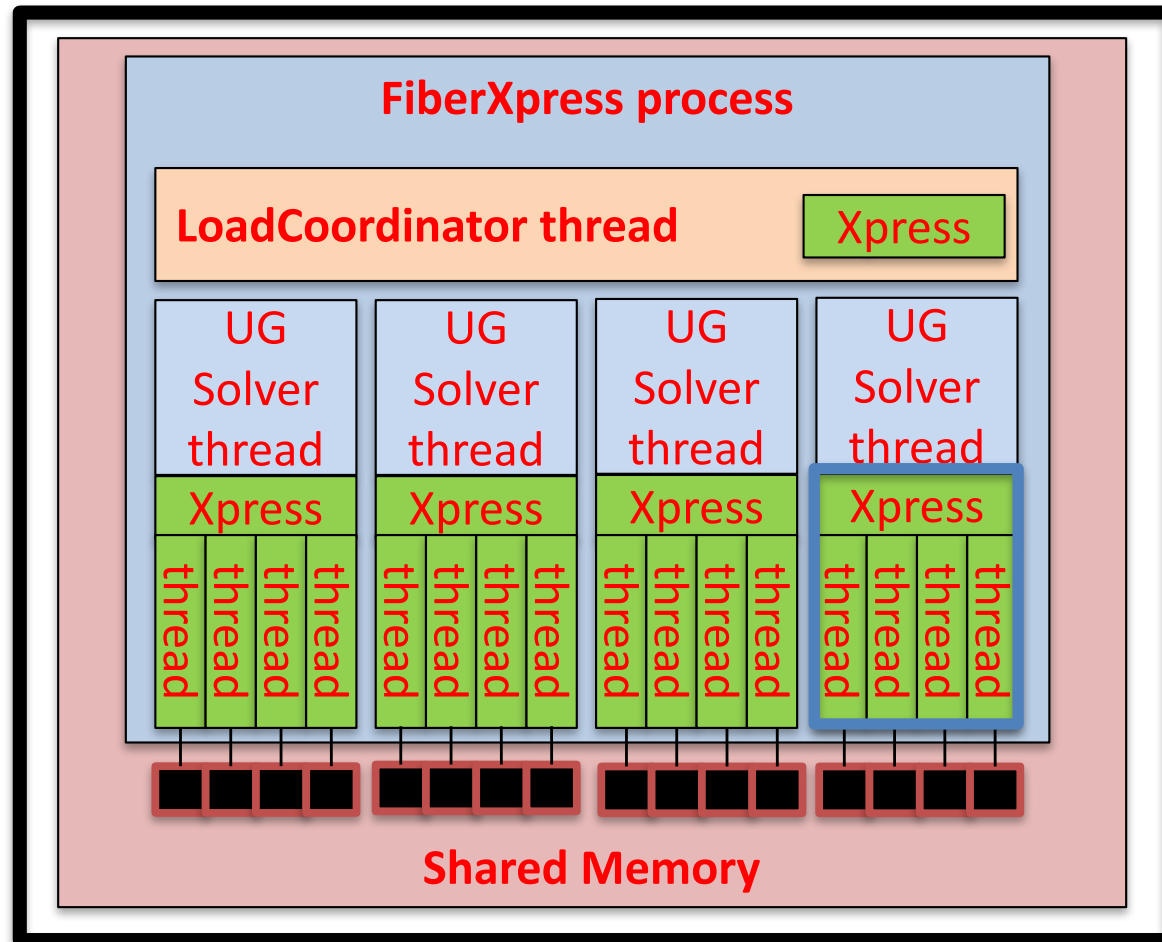
All transfer data need to be converted back for the presolved instance

All feasible solutions need to be converted back for the original instance



# FiberXpress: UG Solver threads with Xpress threads

UG 4 Xpress 4



Berthold T., Farmer J., Heinz S.,  
Perregaard M. (2016)  
Parallelization of the FICO Xpress-  
Optimizer. In: Greuel GM., Koch T.,  
Paule P., Sommesse A. (eds)  
Mathematical Software – ICMS  
2016. ICMS 2016. Lecture Notes in  
Computer Science, vol 9725.  
Springer, Cham

■ : CPU core

Shared Memory Machine (PC)

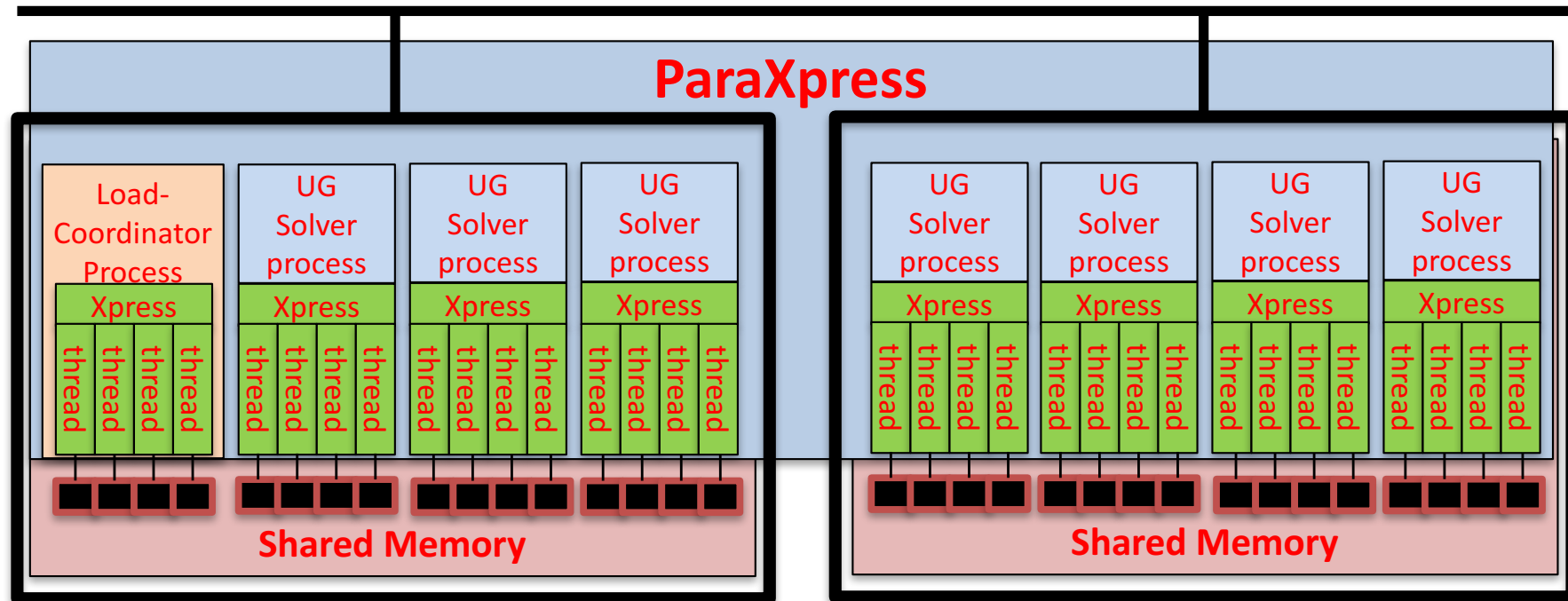


# ParaXpress: UG Solver processes with Xpress threads

## ug[Xpress,MPI] : ParaXpress

- ▶ A powerful massively parallel MIP solver
  - ▶ Can handle, hopefully efficiently, up to  $80,000 \text{ (MPI processes)} \times 24 \text{ (threads)} = 1,920,000 \text{ (cores)}$

Connection network



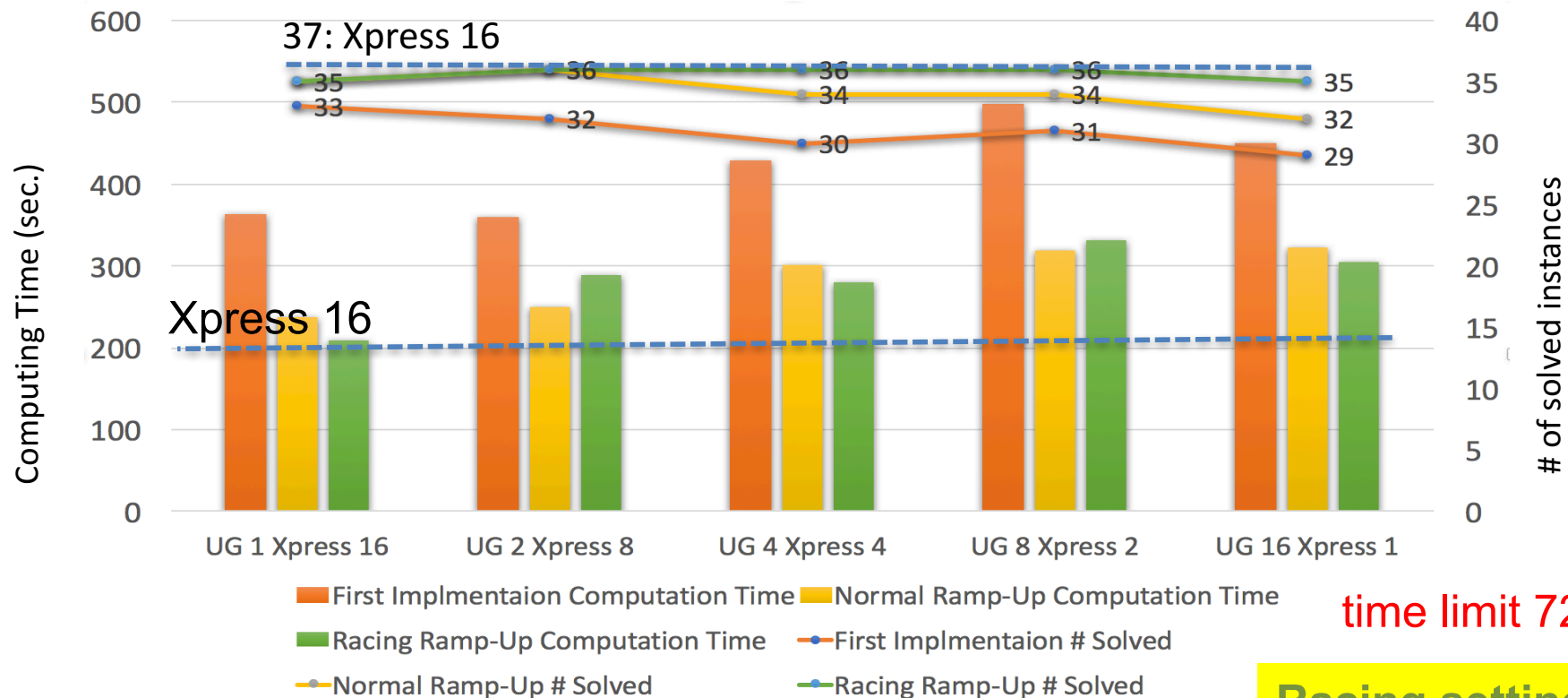
Processing Element or Compute node

Processing Element or Compute node

■ : CPU core



# Computational results of FiberXpress



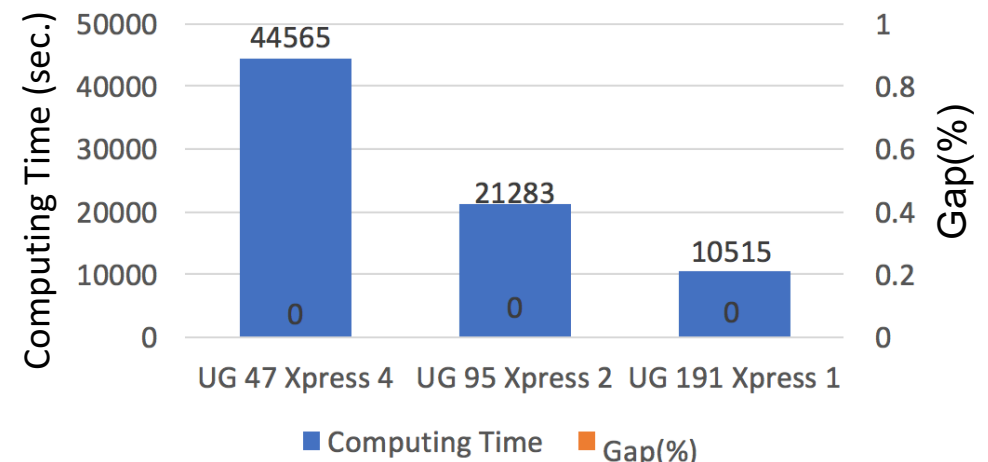
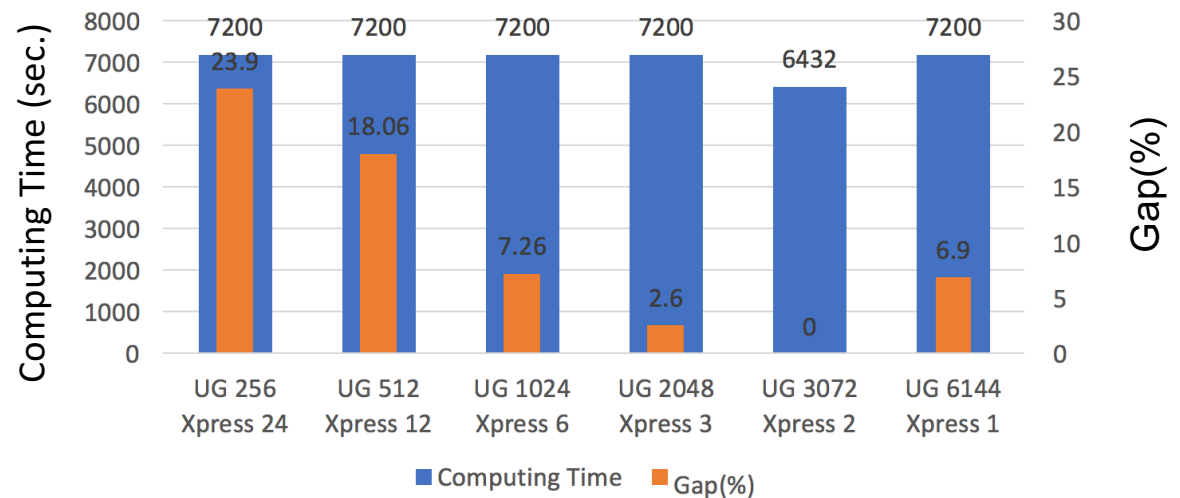
- **Xpress 8.3** and UG-0.8.4 – dev.
- 28 core Intel Xeon CPU E5-2690 v4 CPUs at 2.6 GHz with 35MB cache and 128GB memory
- **52 instances** from Tree test set from MIPLIB2010 (**6** instances were solved at root by Xpress 16)
- Geometric mean (Time out is counted as **7200 sec.**)

Price of externalization is dramatically decreased

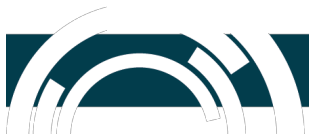
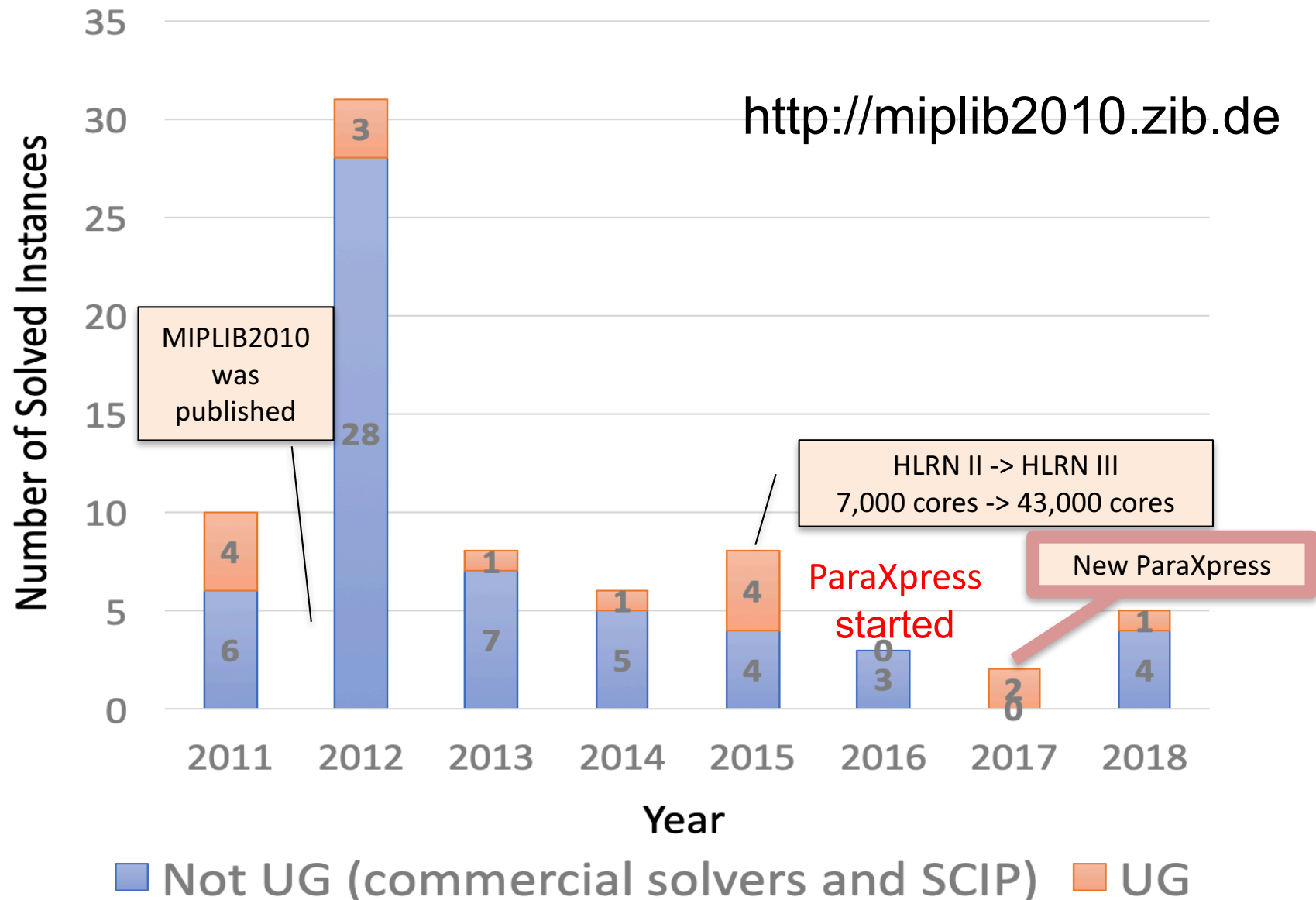


# Computational results of ParaXpress

- Solving timtab2 from MIPLIB2003
  - The **experimental implementation**
    - **Xpress 7.9** and UG-0.8.2
    - HLRN III:
      - Cray XC30 with 24 core Intel E5-2695 v2 CPUs at 2.4GHz, 64GB Memory per computing node
      - 6144 cores are used**
  - New implementation
    - **Xpress 8.3** and UG-0.8.4 – dev.
    - ISM Supercomputer System "I":
      - SGI ICE X with 12 cores Intel Xeon E5-2697v2 at 2.7GHz (2CPUs), 128GB memory per computing node
      - 192 cores are used**



# Solving open instances from MIPLIB2010





# Solving open instances from MIPLIB2017

- Open instances which could not be generated any feasible solution for all solvers

Instance	Solver	#Cores	Time to first sol. (sec.)	Value of first sol.	Time to solve (sec.)	Optimal value
fhnw-sq2	ParaSCIP	72	46,863	0	46,883	0
neos-4409277-trave	ParaSCIP	72	181,150	8	441,575	3
supportcase3	ParaSCIP	72	1,539	0	1,551	0
woodlands09	ParaSCIP	72	218,528	8	245,205	0
neos-3214367-sovi	ParaXpress	72	40,188	180,025	604,810(cont.)	178,326.9644(LB)
	ParaXpress	72	-	179,985(UB)	604,793(cont.)	178,236.7278(LB)
	ParaXpress	2,304	-	179,985(UB)	604,788(cont.)	179,186.6611(LB)
	ParaXpress	2,304	-	179,965(UB)	11,016	179,965
neos-3211096-shag	ParaSCIP	72	-	-	42,484	Infeasible
	ParaXpress	72	-	-	370,983	Infeasible
neos-3631363-vilnia	ParaSCIP	72	-	-	4,632	Infeasible
	ParaXpress	72	-	-	16,335	Infeasible

ISM (Institute of Statistical Mathematics) supercomputer, which is a HPE SGI 8600 with 384 compute nodes, each node has two Intel Xeon Gold 6154 3.0GHz CPUs (18 cores × 2) sharing 384GB of memory, by using ParaSCIP and ParaXpress



# ParaXpress (= ug[Xpress, MPI]) status

---

- ParaXpress can be upgraded new version of Xpress without UG code change
- ParaXpress **is ready to run on over a million CPU cores!**
  - We need an application in which the optimal solution is very important in the research field



# Outline

---

- Notes about the UG design and its dynamic load balancing
- Main computational results of ParaSCIP and ParaXpress
  - Solving previously unsolvable instances of MIP
- Some other projects with ParaSCIP
  - A new feature of SCIP is going to be parallelized
  - An application of `ug[SCIP,*]` libraries
- How UG applications were developed
- Future plan
- Concluding remarks



# A new feature of SCIP (Added by Stephen J. Maher)

## Features

- very fast standalone solver for linear programming (LP), mixed integer programming (MIP), and mixed integer nonlinear programming (MINLP)
- framework for branching, cutting plane separation, propagation, pricing, and Benders' decomposition,
- large C-API, C++ wrapper classes for user plugins

### Parallelization of Benders' decomposition

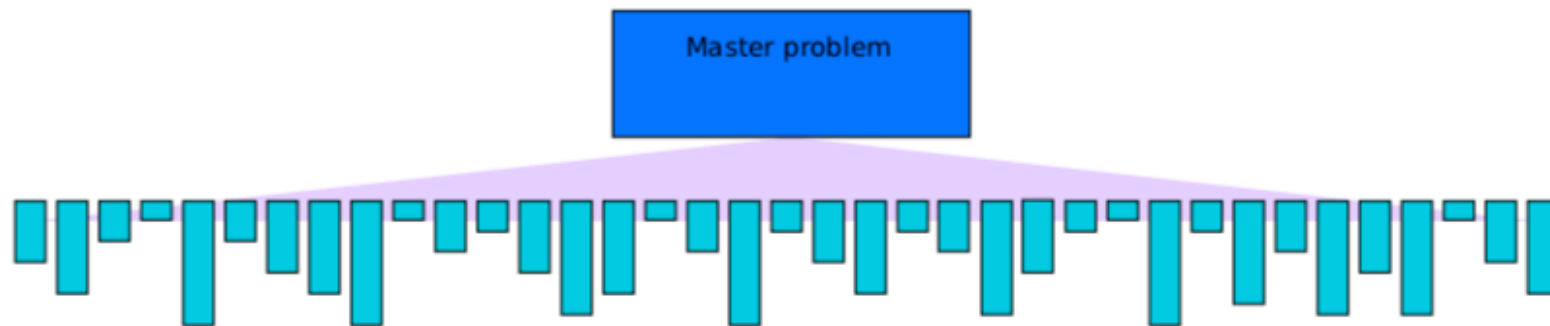
- ▷ Benders' decomposition is easily parallelisable – however, not embarrassingly parallel
- ▷ There are lots of bottlenecks – long running master or subproblems
- ▷ Parallelization of Benders' decomposition can better employ available computational resources
- ▷ For stochastic programs, will help increase the scale of scenarios that can be handled and lead to improved solution quality



# A new feature of SCIP (Added by Stephen J. Maher)

## Three methods of parallelization

1. Tree search parallelization
  - ▶ Using the UG Framework – extended to transfer Benders' cuts between solvers
2. Subproblem parallelization
  - ▶ OpenMP used for shared memory parallelization of subproblem solving
3. Hybrid tree search and subproblem parallelization
  - ▶ Distributed memory tree search parallelization with shared memory subproblem parallelization



# Outline

---

- Notes about the UG design and its dynamic load balancing
- Main computational results of ParaSCIP and ParaXpress
  - Solving previously unsolvable instances of MIP
- Some other projects with ParaSCIP
  - A new feature of SCIP is going to be parallelized
  - An application of `ug[SCIP,*]` libraries
- How UG applications were developed
- Future plan
- Concluding remarks



## Problems

- AIC-based variable selection  
in linear regression and logistic regression  
⇒ mixed integer nonlinear programming problems (MINLP)

## Algorithm proposed in [1, 2]

- Branch-and-bound (B&B) algorithm using **SCIP** and **UG**
- Developed additional plugins:
  - ✦ Relaxation handler (to compute lower bounds efficiently)
  - ✦ Primal heuristics (to find good solutions early)
  - ✦ Branching rules (to reduce branch-and-bound nodes)

- 
- [1] K. Kimura and H. Waki: Minimization of Akaike's information criterion in linear regression analysis via mixed integer nonlinear program. OMS, 33(3), 633–649, 2018.
- [2] K. Kimura: Application of a mixed integer nonlinear programming approach to variable selection in logistic regression. JORSJ, 62(1), to appear.

## Variable Selection

- provides a simple statistical model
- finds a subset of relevant variables
- improves prediction performance

## Direct Objective Optimization in Variable Selection

- $f$ : the goodness-of-fit (how well a model fits a given dataset)
- $\lambda \sum z_j (\lambda > 0)$ : a penalty for the number of variables

- ◆ The problems in [1, 2]  
are of the following form

$$\begin{aligned} \min_{\beta, z} \quad & f(\beta) + \lambda \sum_{j \in J} z_j \\ \text{s.t.} \quad & z_j = 0 \Rightarrow \beta_j = 0 \quad \forall j \in J, \\ & \beta_j \in \mathbb{R}, z_j \in \{0, 1\} \quad \forall j \in J. \end{aligned}$$

- 
- [1] K. Kimura and H. Waki: Minimization of Akaike's information criterion in linear regression analysis via mixed integer nonlinear program. OMS, 33(3), 633–649, 2018.
- [2] K. Kimura: Application of a mixed integer nonlinear programming approach to variable selection in logistic regression. JORSJ, 62(1), to appear.



# Numerical Experiments in [2]

**Problem:** AIC-based variable selection in logistic regression

## Comparison of Approaches:

- Customized B&B algorithm [2]
  - using **SCIP** and **UG**
- Piecewise linear approximation approach [Sato, et al., 2016]
  - using **CPLEX** to solve MILP problems
  - employing existing heuristics for initial solutions

**Benchmark datasets:** UCI machine learning repository

(<https://archive.ics.uci.edu/ml/>)

---

[2] K. Kimura: Application of a mixed integer nonlinear programming approach to variable selection in logistic regression. JORSJ, 62(1), to appear.

# Numerical Experiments in [2]

(16 threads)			Approach 1		Approach 2	
Name	$n$	$p$	AIC	Time(s)	AIC	Time(s)
bumps	2584	22	<b>1097.1</b>	<b>20.0</b>	1098.1	41.5
breast-P	194	34	147.0	<b>25.8</b>	147.0	112.4
biodeg	1055	42	653.3	<b>221.5</b>	653.3	>5000
stat-G	1000	62	958.2	>5000	958.2	>5000
madelon	2000	500	<b>2502.1</b>	>5000	2504.0	>5000

- Approach 1: Customized B&B algorithm [2] (**SCIP** and **UG**)
- Approach 2: Piecewise linear approximation approach (**CPLEX**)
- $n$ : the number of observations
- $p$ : the number of candidates for variables

# Outline

---

- Notes about the UG design and its dynamic load balancing
- Main computational results of ParaSCIP and ParaXpress
  - Solving previously unsolvable instances of MIP
- Some other projects with ParaSCIP
  - A new feature of SCIP is going to be parallelized
  - An application of `ug[SCIP,*]` libraries
- How UG applications were developed
- Future plan
- Concluding remarks



# Instantiated parallel solvers by UG

---

The following solvers have developed by Yuji Shinano  
cooperated with each solver developers

- Single thread base solver
  - ParaSCIP: `ug[SCIP, MPI]`, FiberSCIP: `ug[SCIP, Pthreads/C++11]`
- Multi-threaded base solver
  - ParaXpress: `ug[Xpress, MPI]`, FiberXpress: `ug[Xpress, Pthreads/C++11]`
- Concorde: TSP solver
  - `ug[Concorde, MPI]`, `ug[Concorde, Pthreads/C++11]`
  - `ug[Concorde/MPI, MPI]`
    - **Concorde/MPI** has developed by Utz-Uwe Haus



# Instantiated parallel solvers by UG

---
































The following solvers have developed  
by each base solver developers.

Skelton code was distributed.  
Design and coding together with Yuji Shinano  
about one week in the beginning

- Multi-threaded base solver
  - ParaNUOPT: `ug[NUOPT, MPI]`, FiberNUOPT: `ug[NUOPT, Pthreads/C++11]`
    - Only tested with single thread NUOPT
- Distributed base solver
  - `ug[PIPS-SBB, MPI]`
    - PIPS-SBB: a distributed memory solver for Stochastic MIP



# Skelton code

 baseSolverParaComm.h	 baseSolverParaInstance.h
 baseSolverParaCommMpi.cpp	 baseSolverParaInstanceMpi.cpp
 baseSolverParaCommMpi.h	 baseSolverParaInstanceMpi.h
 <b>baseSolverParaCommPointHdlr.cpp</b>	 baseSolverParaInstanceTh.cpp
 baseSolverParaCommPointHdlr.h	 baseSolverParaInstanceTh.h
 baseSolverParaCommTh.cpp	 baseSolverParaSolution.cpp
 baseSolverParaCommTh.h	 baseSolverParaSolution.h
 baseSolverParaDef.h	 baseSolverParaSolutionMpi.cpp
 baseSolverParaDeterministicTimer.h	 baseSolverParaSolutionMpi.h
 baseSolverParaDiffSubproblem.cpp	 baseSolverParaSolutionTh.cpp
 baseSolverParaDiffSubproblem.h	 baseSolverParaSolutionTh.h
 baseSolverParaDiffSubproblemMpi.cpp	 baseSolverParaSolver.cpp
 baseSolverParaDiffSubproblemMpi.h	 baseSolverParaSolver.h
 baseSolverParaDiffSubproblemTh.h	 fbaseSolver.cpp
 baseSolverParaInitiator.cpp	 parabaseSolver.cpp
 baseSolverParaInitiator.h	

Rename: baseSolver->your\_solver\_name, Fill codes for all virtual functions  
baseSolverPara**CommPoint**Hdlr.cpp: **CommPoint** is a callback of solver  
to communicate with LoadCorrdinator



# Built parallel solvers with ug[SCIP,\*] libraries

---

The following solvers have developed  
by each SCIP plugins developers.

Debugged with Yuji Shinano

- ug[SCIP-Jack,\*]
  - Parallel solvers for solving Steiner Tree Problems
- ug[SCIP-SDP,\*]
  - Parallel solvers for solving Mixed Integer Semidefinite Programs



# Outline

---

- Notes about the UG design and its dynamic load balancing
- Main computational results of ParaSCIP and ParaXpress
  - Solving previously unsolvable instances of MIP
- Some other projects with ParaSCIP
  - A new feature of SCIP is going to be parallelized
  - An application of ug[SCIP,\*] libraries
- How UG applications were developed
- Future plan
- Concluding remarks





# Plan for UG 1.0

---

- If you have your own state-of-the-art solver, please try to parallelize it with UG
- If you have SCIP application, please try to parallelize it with `ug[SCIP,*]` libraries

We hope to have feedbacks from user side

- Second UG workshop will be held in two years, since “Sustainable Infrastructures for Archiving and Publishing High-Performance Optimization Software” is three years project
  - Right after the workshop, UG 1.0 would be released



# Outline

---

- Notes about the UG design and its dynamic load balancing
- Main computational results of ParaSCIP and ParaXpress
  - Solving previously unsolvable instances of MIP
- Some other projects with ParaSCIP
  - A new feature of SCIP is going to be parallelized
  - An application of ug[SCIP,\*] libraries
- How UG applications were developed
- Future plan
- Concluding remarks



# Concluding remarks

---

- A research platform for parallel branch-and-bound
  - Users can run all parallel solvers by themselves for fair direct comparison
    - source codes are publically available
    - test data are publically available

